Homework 5: Divide & Conquer

Due: October 17, 2024

Problem 1. The recurrence $T(n) = 7T(n/2) + n^2$ describes the running time of an algorithm A. A competing algorithm A' has a running time of $T'(n) = aT'(n/4) + n^2$. What is the largest integer value for a such that A' is asymptotically faster than A?

Solution.

Problem 2. The Euclidean algorithm is a method for computing the greatest common divisor (gcd) of two numbers, taking advantage of the fact

$$gcd(a, b) = gcd(b, a - b)$$

for positive integers a, b satisfying $a \ge b$. Consider the following variation:

$$gcd(a,b) = \begin{cases} 2 gcd(a/2, b/2) & \text{if } a, b \text{ even} \\ gcd(a, b/2) & \text{if } a \text{ odd } b \text{ even} \\ gcd((a-b)/2, b) & \text{if } a, b \text{ odd} \end{cases}$$

- a. Prove the variation is correct.
- b. Provide an algorithm that uses the variation to compute the greatest common divisor of two numbers a, b in $O(\log(ab))$

Solution.

Problem 3. Given an n-bit binary integer, design a divide-and-conquer algorithm to convert it into its decimal representation. For simplicity, you may assume that n is a power of 2.

- 1. Provide a succinct (but clear) description of your algorithm, including pseudocode.
- 2. Prove the correctness of your algorithm.
- 3. Analyze the running time of your algorithm. Assume that it is possible to multiply two decimal integers numbers with at most m digits in $O(m^{\log_2 3})$ time.

Hint: An *n*-bit binary integer *x* can be expressed as $x = (x_{n-1}, x_{n-2}, \dots, x_1, x_0)_2$ where $x_i \in \{0, 1\}$. Let $x_{\ell} = (x_{n/2-1}, x_{n/2-2}, \dots, x_1, x_0)_2$ be the (n/2)-bit binary integer corresponding to the (n/2) least significant digits of *x*. Let $x_m = (x_{n-1}, x_{n-2}, \dots, x_{n/2+1}, x_{n/2})_2$ be the (n/2)-bit binary integer representing the (n/2) most significant digits of *x*. Then, $x = x_{\ell} + 2^{n/2} \cdot x_m$. This should suggest us a way to set up a divide and conquer strategy...:) Careful about the number of subproblems!

Solution.